# A DIRECT ALGORITHM FOR DISTINGUISHING NONSINGULAR $M$－MATRIX AND $H$－MATRIX＊ 

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#### Abstract

A\) direct algorithm is proposed by which one can distinguish whether a ma－ trix is an M－matrix（or H－matrix）or not quickly and effectively．Numerical examples show that it is effective and convincible to distinguish M－matrix（or H－matrix）by using the algorithm．


Key words nonsingular M－matrix，nonsingular H－matrix，direct algorithm．
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## 1 Introduction

For many kinds of applications of $M$－matrices and $H$－matrices，the problem how to deter－ mine whether a matrix is an $M$－matrix（or $H$－matrix）or not arouses many researchers interesting． Recently，some iterative methods have been proposed for distinguishing $H$－matrices（see［2－5］）． However，these methods have a common drawback，that is，it is not possible to determine the number of steps of iteration，and when $A$ is not an $H$－matrix，a wasteful computation is necessary． A direct algorithm has been proposed in［6］，but it is only useful when matrices are symmetrical． In this paper，to conquer these drawbacks，we propose a new direct algorithm．

## 2 A direct algorithm for distinguishing $M$－matrix

Let $R^{n \times n}$ denote the set of all $n \times n$ real matrices．$A=\left(a_{i j}\right) \in R^{n \times n}$ is said to be an $M$－matrix if $a_{i j} \leq 0$ ，for $i \neq j$ ，and $A^{-1} \geq 0$ ．

Lemma $1^{[1]}$ Let $A=\left(a_{i j}\right) \in R^{n \times n}$ be an $M$－matrix，then any principle submatrix of $A$ is an $M$－matrix．

Lemma $2^{[1]}$ Let $A=\left(a_{i j}\right) \in R^{n \times n}$ ，its off－diagonal entries are all non－positive，then $A$ is

[^0]an $M$-matrix if and only if successive principle minor of $A, D_{K}>0, k=1, \cdots, n$.
From Lemma 2, we can immediately obtain the following lemma.
Lemma 3 Let $A=\left(a_{i j}\right) \in R^{2 \times 2}$, and $a_{i j} \leq 0, i \neq j, a_{i i}>0$, then $A$ is an $M$-matrix if and only if determinant of $A, \operatorname{det} A>0$.

Theorem 1 Let

$$
B=\left[\begin{array}{ll}
B_{11} & B_{12} \\
B_{21} & B_{22}
\end{array}\right] \in R^{n \times n},
$$

where $B_{12} \leq 0, B_{21} \leq 0, B_{11}$ is a $2 \times 2$ square matrix and $B_{22}$ is an $(n-2) \times(n-2)$ square matrix, in which their diagonal entries are all positive and off-diagonal entries are all non-positive. Then $B$ is an $M$-matrix if and only if $\operatorname{det} B_{11}>0$ and $B_{22}-B_{21} B_{11}^{-1} B_{12}$ is an $M$-matrix.

Proof Necessity: Suppose $B$ is an $M$-matrix, then
$B^{-1}=\left[\begin{array}{cc}B_{11}^{-1}+B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} & -B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} \\ -\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} & \left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1}\end{array}\right] \geq 0$, and $B_{11}$ and $B_{22}$ are $M$-matrices by Lemma 1. Hence, $\operatorname{det} B_{11}>0$ by Lemma 3 , and $B_{11}^{-1} \geq$ $0,\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} \geq 0$. For $B_{12} \leq 0, B_{21} \leq 0$, we have $B_{21} B_{11}^{-1} B_{12} \geq 0$, and off-diagonal entries of matrix $B_{22}-B_{21} B_{11}^{-1} B_{12}$ are all non-positive. So, $B_{22}-B_{2} B_{11}^{-1} B_{12}$ is an $M$-matrix.

Sufficiency: Suppose $\operatorname{det} B_{11}>0$ and $B_{22}-B_{21} B_{11}^{-1} B_{12}$ is an $M$-matrix, then by Lemma 3, we have that $B_{11}$ is an $M$-matrix, so

$$
\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} \geq 0, \quad B_{11}^{-1} \geq 0 .
$$

Therefore

$$
\begin{aligned}
& B_{11}^{-1}+B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} \geq 0, \\
& -\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} \geq 0, \\
& -B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} \geq 0 .
\end{aligned}
$$

From these inequalities, we have
$B^{-1}=\left[\begin{array}{cc}B_{11}^{-1}+B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} & -B_{11}^{-1} B_{12}\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} \\ -\left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1} B_{21} B_{11}^{-1} & \left(B_{22}-B_{21} B_{11}^{-1} B_{12}\right)^{-1}\end{array}\right] \geq 0$.
Thus $B$ is an $M$-matrix.
From Theorem 1, we propose the following algorithm A.

## Algorithm A

Input The given matrix $B=\left(b_{i j}\right) \in R^{n \times n}$.
Step 1 Set $B=B^{(m)}$, and $m=0$.
Step 2 Partition $B^{(m)}$ into a $2 \times 2$ block matrix

$$
B^{(m)}=\left[\begin{array}{ll}
B_{11}^{(m)} & B_{12}^{(m)} \\
B_{21}^{(m)} & B_{22}^{(m)}
\end{array}\right],
$$


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