

## MATRIX ANALYSIS ON LEADING TERM OF CONDITION NUMBER FOR ADDITIVE SCHWARZ METHODS<sup>\*1)</sup>

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### Abstract

It is well known the order of preconditioned matrix by using additive Schwarz methods. In order to estimate the resulted PCG iteration counts, the related leading term before the order is given in this paper.

*Key words:* Additive Schwarz method, PCG, Matrix analysis.

### 1. Introduction

Let us consider the following second order elliptic boundary value problem:

$$\mathcal{L}u = f \quad \text{in } \Omega \quad (1)$$

$$u = 0 \quad \text{on } \partial\Omega \quad (2)$$

where  $\mathcal{L}$  is a self-adjoint positive operator and

$$\Omega \subset \mathbb{R}^d \quad (1 \leq d \leq 3)$$

is a polyhedral domain.

Using weak solution it leads to a discrete equation

$$Au = f \quad (3)$$

with

$$A = (\alpha_{ij}), \quad \alpha_{ij} = \mathcal{A}(\phi_i, \phi_j) \quad (4)$$

where  $\{\phi_i\}$  could be nodal basis consisting of piece-wise linear functions or other spline functions. It is well known that the coefficient matrix  $A$  is symmetry positive definite matrix with condition number

$$\kappa(A) := \frac{\lambda_{max}(A)}{\lambda_{min}(A)} = O(h^{-2}) \quad (5)$$

Furthermore, under rectangle uniform mesh it is easy to obtain the leading term of Discrete Laplacian condition number before the above order as follows which is independent on dimension.

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Dimension	Scheme	$\lambda_{max}$	$\lambda_{min}$	$\kappa$
1 - D	3 - point	$2 - (\pi h)^2$	$(\pi h)^2$	$\frac{2}{\pi^2} h^{-2}$
2 - D	5 - point	$4 - 2(\pi h)^2$	$2(\pi h)^2$	$\frac{2}{\pi^2} h^{-2}$
3 - D	7 - point	$6 - 3(\pi h)^2$	$3(\pi h)^2$	$\frac{2}{\pi^2} h^{-2}$

When we use conjugate gradient algorithm for solving the system, for a given tolerance, the iteration number will proportional to  $h^{-1}$ . This convergent rate is really slow for a large scale problems. It is our aim to study how better is better the preconditioner by using additive Schwarz method.

Suppose there are two subdomain partitions, one is without overlapping

$$\Omega = \cup \Omega_i \text{ with } \Omega_i \cap \Omega_j = \emptyset \text{ if } i \neq j$$

and another is with overlapping

$$\Omega = \cup \hat{\Omega}_i \text{ with } \hat{\Omega}_i \cap \hat{\Omega}_j \neq \emptyset \text{ if } i \neq j$$

and denote  $A_i$  be a matrix representation which is the restriction of the original operator  $A$  over the subdomain  $\Omega_i$ . Without considering unknowns permutation appropriately, then the related preconditioner can be written as

$$B_i = \begin{bmatrix} A_i^{-1} & 0 \\ 0 & 0 \end{bmatrix}$$

and the corresponding matrix  $A$  has the following block form

$$A = \begin{bmatrix} A_i & A_{ij} \\ A_{ij}^t & A_j \end{bmatrix}, \quad B_i A = \begin{bmatrix} I & A_i^{-1} A_{ij} \\ 0 & 0 \end{bmatrix}$$

Hence, in general we obtain a whole subdomain preconditioner  $B_s$  as follows:

$$B_s = \sum B_i, \text{ with } B_i = R_i^t A_i^{-1} R_i \quad (6)$$

where  $R_i$  is a truncated permutation matrix from  $\Omega$  to the subdomain  $\Omega_i$ .

Some questions now are arisen on the above algorithm as a PCG preconditioner . What is efficiency of the above parallel ASM ? How fast the related PCG iteration does converge? How does the eigenvalues of preconditioned matrix  $B_s A$  distribute ? How to estimate the largest and smallest eigenvalues of  $B_s A$  ?

This estimation can be reduced to special generalized eigen-decomposition

$$Au = \lambda B_s^{-1} u$$

From finite element theory, it has been known by using energy norm estimation

$$\lambda_{min}(B_s A) = O((hH)^{-1}), \quad \lambda_{max}(B_s A) = O(1).$$

where  $h$  is the smallest mesh size on subdomains, and  $H$  is the subdomain width.

In practical, it is need to estimation the leading term before the order to estimate the PCG iteration counts more precisely for a given tolerance.

## 2. Some Properties of Related Matrices $Q_s$ and $Q_\Gamma$

With our assumption,  $A$  is a M-matrix. Hence, matrix each  $B_i$  is non-negative and  $B_i A$  is a matrix with all non-positive off-diagonals, so does  $B_s A$ . Based on discrete maximum principle, it is not hard to show that

$$\lambda_{\max}(B_s A) \leq 1 + \max\{d_i\},$$

where  $d_i$  – overlapping multiple index among subdomains.

In order to estimate the smallest eigenvalue of preconditioned matrix  $B_s A$  more clearly, we turn to study a related matrix  $Q_s = I - B_s A$ . It is clear that after a proper permutation  $Q_s$  can be represented in terms of

$$Q_s = \begin{bmatrix} D & * \\ 0 & Q_\Gamma \end{bmatrix}$$

where  $D$  is diagonal with non-positive integer

$$D = \text{diag}\{1 - d_i\},$$

We call  $Q_\Gamma$  as equivalent interior boundary matrix to  $Q_s$ . It is clear that  $Q_\Gamma$  consists from  $Q_s$  deleting related columns with zero on all non-diagonals and related rows. The size of  $Q_\Gamma$  is much smaller than  $Q_s$ . In fact  $Q_\Gamma$  can be seen as a restriction on the discrete overlapping region of  $Q_s$ . All eigenvalues of  $Q_\Gamma$  belong to  $\sigma(Q_s)$ . In particular, if  $\lambda_{\max}$  is the largest eigenvalue of  $Q_\Gamma$ , then  $\mu_{\min} = 1 - \lambda_{\max}$  is the smallest eigenvalue of  $B_s A$ .

Hence, the smallest eigenvalue of  $B_s A$  is just the same as the largest eigenvalue of  $Q_\Gamma$  which is only defined on the much smaller discrete overlapping region. Furthermore, it is easy to know  $Q_\Gamma$  is non-negative with all zeroes along the main diagonals. The eigenvector  $u_\Gamma$ , corresponding  $\lambda_{\max}$ , is non-negative, too. In another word, the matrix  $Q_\Gamma$  has unique positive eigenvector  $u_\Gamma$ , and  $Q_\Gamma u_\Gamma = \lambda_{\max} u_\Gamma$ . In the matrix point of view, to get an estimation of  $\lambda_{\max}$   $\mu \leq \lambda_{\max} \leq \lambda$ , it is sufficient to find a positive vector  $z$  such that  $\mu z \leq Q_\Gamma z \leq \lambda z$ ,

## 3. 1-D Case: Model Problem

$$-(pu')' = f, \quad u(0) = u(1) = 0 \tag{7}$$

where  $p(x) \geq p_0 > 0$ . For a given mesh partition

$$\Delta : 0 = x_0 < x_1 < \dots < x_{N+1} = 1$$

with  $\delta$  mesh size overlapping subdomain partition

$$s_1 = n_1, s_2 = n_1 + n_2, s_m = \sum_{k=1}^m n_k = N, \quad \text{with } n_k > 1, k = 1, 2, \dots, m$$

$$H_1 = x_{s_1+\delta} - x_0, \quad H_k = x_{s_k+\delta} - x_{s_{k-1}-\delta}, \quad (s = 2, \dots, m-1) \quad H_m = x_1 - x_{s_{m-1}-\delta}.$$

By using finite element method or difference method, the resulting stiffness matrix  $A$  is tridiagonal of order  $N$  and its coefficients are equal to

$$\alpha_{ij} = \mathcal{A}(\phi_i, \phi_j) = \begin{cases} 0 & \text{if } |i-j| > 1 \\ -p_{i-1/2} h_i^{-1} & \text{if } j = i-1 \\ p_{i-1/2} h_i^{-1} + p_{i+1/2} h_{i+1}^{-1} & \text{if } j = i \\ -p_{i+1/2} h_{i+1}^{-1} & \text{if } j = i+1 \end{cases} \tag{8}$$

where  $p_{i-1/2} = (p_{i-1} + p_i)/2$ ,  $p_{i+1/2} = (p_{i+1} + p_i)/2$ .

If we introduce so-called generalized mesh  $\hat{h}_i = \frac{h_i}{p_{i-1/2}}$ , the general variable coefficient case can be taken as a non-uniform mesh with constant coefficients.

Using ASM to the above matrix  $A$  leads to

$$B_k A = \begin{bmatrix} 0 & 0 & & \dots \\ \dots & \dots & & \dots \\ \dots & A_{k,k}^{-1} A_{k,k-1} & I & A_{k,k}^{-1} A_{k,k+1} \\ 0 & 0 & & \dots \\ \dots & \dots & & \dots \end{bmatrix},$$

where  $\text{rank}(A_{11}) = n_1 + \delta$ ,  $\text{rank}(A_{kk}) = n_k + 2\delta$ ,  $k = 2, \dots, m-1$ ,  $\text{rank}(A_{mm}) = n_m + \delta$ .

Let  $A_{k,k}^{-1} A_{k,k-1} = \{0, 0, \dots, 0, u^{k-}\}$ ,  $A_{k,k}^{-1} A_{k,k+1} = \{u^{k+}, 0, \dots, 0, 0\}$ . where  $u^{k-}$  and  $u^{k+}$  are the solution of related two-point boundary value problems as follows

$$\mathcal{L}u = 0, \quad u(x_{n_k-\delta}) = 1, \quad u(x_{n_k+\delta}) = 0;$$

and

$$\mathcal{L}u = 0, \quad u(x_{n_k+\delta}) = 1, \quad u(x_{n_k-\delta}) = 0;$$

The interior boundary matrix  $Q_\Gamma = (I - B_s A)|_\Gamma$  becomes

$$Q_\Gamma = \begin{bmatrix} 0_2 & q_{12} & 0_2 & \dots \\ q_{21} & 0_2 & q_{23} & 0_2 & \dots \\ \dots & & \dots & & \dots \\ \dots & q_{j,j-1} & 0_2 & q_{j,j+1} & \dots \\ \dots & & \dots & & \dots \\ \dots & & 0_2 & q_{m+1,m} & 0_2 \end{bmatrix},$$

where  $0_2 - 2 \times 2$  zero matrix, caused from two-point boundary value problem.

$$0_2 = \begin{bmatrix} 0 & 0 \\ 0 & 0 \end{bmatrix}, \quad q_{j,j-1} = \begin{bmatrix} 0 & -(1 - \beta_{j-1}) \\ 0 & -\alpha_j \end{bmatrix}, \quad q_{j,j+1} = \begin{bmatrix} -\beta_{j-1} & 0 \\ -(1 - \alpha_j) & 0 \end{bmatrix}$$

In the uniform mesh case  $\alpha_j = \alpha$ ,  $\beta_j = \beta$

$$\Delta_k(\lambda) = \{\lambda^2 - (1 - \alpha - \beta)\} \Delta_{k-1}(\lambda) - \lambda^2 \alpha \beta \Delta_{k-2}(\lambda),$$

$$\Delta_0(\lambda) = 1, \quad \Delta_1(\lambda) = \lambda^2 - (1 - \alpha)(1 - \beta).$$

A general representation of  $\Delta_k(\lambda)$  was given by Jiachang Sun and Tony Chan in [1] as follows

$$\Delta_m(\lambda) = \frac{(\alpha\beta)^{m/2} \lambda^{m-1}}{\sin \phi} \{\lambda \sin(m+1)\phi - (\alpha\beta)^{1/2} \sin m\phi\}, \quad (9)$$

where  $\cos \phi = \frac{\lambda^2 - (1 - \alpha - \beta)}{2(\alpha\beta)^{1/2}\lambda}$ . Thus, there is an estimation of eigenvalues for small  $\alpha, \beta$

$$\lambda_k = (\alpha\beta)^{1/2} \cos \phi_k \pm \sqrt{\alpha\beta(\cos \phi_k)^2 + (1 - \alpha - \beta)}, \quad (10)$$

where  $\phi_k = \frac{k\pi}{m+1}$ , ( $k = 1, 2, \dots, m$ ).

In particular, the minimum eigenvalue is equal to [1]

$$\lambda_{\min}(B_s A) \approx \frac{\pi^2}{2} H h.$$

For non-uniform case, it is not difficult to show that the above argument still holds if  $H$  and  $h$  are replaced by the corresponding average mesh width of subdomain  $H_o$  and minimal generalized mesh width  $\hat{h}_o$  of overlapping among subdomains, respectively:

$$\lambda_{\max}(Q_s) = \lambda_{\max}(Q_\Gamma) \approx 1 - \frac{\pi^2}{2} H_o \hat{h}_o, \quad (11)$$

It is worth to point that eigenvalues of  $Q_s$  depend on partition of subdomains in the overlapping area only. They are independent upon partition within non-overlapping of any subdomains.

#### 4. Main Idea for Reducing Dimension Procedure

In this section we describe the main idea how to estimate a high-dimension eigenvalue problem via related lower dimension case, if the original domain can decomposed by tensor product of related lower dimension.

**Step 1:** Find related lower dimension interior boundary eigenvalue.

**Step 2:** Extended the interior eigenvector to whole lower dimension by solving piecewise discrete boundary value problems on each subdomain with overlapping.

**Step 3:** Expand the lower dimension eigenvector to the higher dimension eigenvector via tensor product.

**Step 4:** Localization required eigenvalue on original dimension by using comparison theorem based on maximum principle.

There are some special cases as follows:

**Case 1.** Reducing dimension degree: from 2-D to 1-D .

Suppose a 2-D discrete operator can be decomposed into 1-D

$$\mathcal{L} = \mathcal{L}_x + \mathcal{L}_y$$

and the discrete domain can also be decomposed as

$$\Omega^h = \Omega_x^h \otimes \Omega_y^h$$

. In this case we can decompose the discrete matrix via Kronecker product

$$A = A_x \otimes I_y + I_x \otimes A_y$$

and we may take an overlapping preconditioner as

$$B_s^{-1} = B_{sx}^{-1} \otimes I_y + I_x \otimes B_{sy}^{-1}.$$

**Case 2.** Reducing dimension degree: from 3-D to 1-D directly.

Suppose a 3-D discrete operator can be decomposed into 1-D

$$\mathcal{L} = \mathcal{L}_x + \mathcal{L}_y + \mathcal{L}_z.$$

and the discrete domain can also be decomposed as

$$\Omega^h = \Omega_x^h \otimes \Omega_y^h \otimes \Omega_z^h.$$

In this case we can decompose the discrete matrix via Kronecker product

$$A = A_x \otimes (I_y \otimes I_z) + I_x \otimes A_y \otimes I_z + (I_x \otimes I_y) \otimes I_z.$$

and we may take an overlapping preconditioner as

$$B_s = B_{sx}^{-1} \otimes (I_y \otimes I_z) + I_x \otimes B_{sy}^{-1} \otimes I_z + (I_x \otimes I_y) \otimes B_{sz}^{-1}.$$

**Case 3.** Reducing dimension degree: from 3-D to 2-D and 1-D.

Suppose a 3-D discrete operator can be decomposed into 2-D plus 1-D

$$\mathcal{L} = \mathcal{L}_{x,y} + \mathcal{L}_z.$$

and the discrete domain can also be decomposed as

$$\Omega^h = \Omega_{x,y}^h \otimes \Omega_z^h.$$

In this case we can decompose the discrete matrix via Kronecker product

$$A = A_{x,y} \otimes I_z + I_{x,y} \otimes A_z.$$

and we may take an overlapping preconditioner as

$$B_s = B_{sx,y}^{-1} \otimes I_z + I_{x,y} \otimes B_{sz}^{-1}$$

## 5. Some Examples with First Boundary Value Condition

We consider rectangle domain in the following examples.

**Example 1.**

$$\mathcal{L}u = -\frac{\partial}{\partial x}a_1(x)\frac{\partial u(x,y)}{\partial x} - \frac{\partial}{\partial y}a_2(y)\frac{\partial u(x,y)}{\partial y}, \quad (12)$$

$$\begin{aligned} u|_{\partial\Omega} &= f \\ \mathcal{L}_x u &= -\frac{\partial}{\partial x}a_1(x)\frac{\partial u(x,y)}{\partial x}, \\ \mathcal{L}_y u &= -\frac{\partial}{\partial y}a_2(y)\frac{\partial u(x,y)}{\partial y} \end{aligned}$$

where  $a_1(x) > 0, a_2(y) > 0$ .

Suppose the discrete domain can also be decomposed as

$$\Omega^h = \Omega_x^h \otimes \Omega_y^h.$$

Thus, via Kronecker product we can decompose the discrete  $M-$  matrix as

$$A = A_x \otimes I_y + I_x \otimes A_y$$

and the related overlapping preconditioner as

$$B_s^{-1} = B_{sx}^{-1} \otimes I_y + I_x \otimes B_{sy}^{-1}.$$

**Step 1:** Find related 1-D interior boundary largest eigenvalue.

$$Q_{\Gamma x} u_{\Gamma x} = \lambda_x u_{\Gamma x}, \quad u_{\Gamma x} \geq 0,$$

$$Q_{\Gamma y} u_{\Gamma y} = \lambda_y u_{\Gamma y}, \quad u_{\Gamma y} \geq 0,$$

where both eigenvectors  $u_{\Gamma x}$  and  $u_{\Gamma y}$ , corresponding to own largest eigenvalue of each positive matrix  $Q_{\Gamma x}$  and  $Q_{\Gamma y}$ , are positive.

**Step 2:** Extended the interior eigenvector to whole 1-D via piecewise discrete boundary value problems.

$$\begin{aligned}\mathcal{L}_x u_x &= 0, \quad \text{in } \Omega_{x,i}^h, \\ u_x &= u_{\Gamma x} \quad \text{on } \partial\Omega_{x,i}^h\end{aligned}$$

and

$$\begin{aligned}\mathcal{L}_y u_y &= 0, \quad \text{in } \Omega_{y,j}^h, \\ u_y &= u_{\Gamma y} \quad \text{on } \partial\Omega_{y,j}^h\end{aligned}$$

Note that, under the assumption, the above two vectors  $u_x$  and  $u_y$  keep positive based on maximum principle. And vectors  $B_{sx}^{-1}u_x$  and  $B_{sy}^{-1}u_y$  are all non-negative.

**Step 3:** Expand the 1-D vector to a 2-D vector via Kronecker product. Set a 2-D positive vector  $u$  such that  $u = u_x \otimes u_y$

Then

$$\begin{aligned}Au &= (A_x \otimes I_y + I_x \otimes A_y)(u_x \otimes u_y) = (A_x u_x) \otimes u_y + u_x \otimes (A_y u_y) \\ &= \lambda_x (B_{sx}^{-1}u_x) \otimes u_y + \lambda_y u_x \otimes (B_{sy}^{-1}u_y)\end{aligned}$$

Note that all values and vectors appeared in the above formula are non-negative.

**Step 4:** Localization of the original smallest eigenvalue.

$$\min(\lambda_x, \lambda_y) B_s^{-1} u \leq Au \leq \max(\lambda_x, \lambda_y) B_s^{-1} u$$

Finally, it leads to an estimation of the smallest eigenvalue for  $B_s A$  as follows

$$\min(\lambda_x, \lambda_y) \leq \lambda_{\min}(B_s A) \leq \max(\lambda_x, \lambda_y)$$

**Example 2.** Consider general 2-D second order self-adjoint case

$$\mathcal{L}u = -\frac{\partial}{\partial x} a_1(x, y) \frac{\partial u(x, y)}{\partial x} - \frac{\partial}{\partial y} a_2(x, y) \frac{\partial u(x, y)}{\partial y} \quad (13)$$

$$\mathcal{L}_x u = -\frac{\partial}{\partial x} a_1(x, y) \frac{\partial u(x, y)}{\partial x}, \quad \mathcal{L}_y u = -\frac{\partial}{\partial y} a_2(x, y) \frac{\partial u(x, y)}{\partial y}$$

where  $a_1(x, y) \geq a_{10} > 0, a_2(x, y) \geq a_{20} > 0$ .

Let

$$\lambda_1 = \min_{y \in \Gamma} \{\lambda_{\min}(B_1 A_1(y))\},$$

$$\lambda_2 = \min_{x \in \Gamma} \{\lambda_{\min}(B_2 A_2(x))\}$$

then

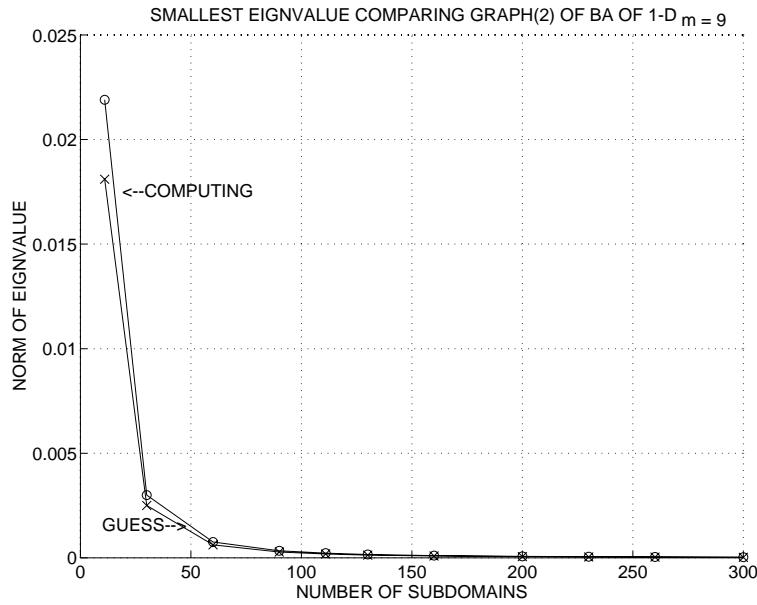
$$\min\{\lambda_1, \lambda_2\} \leq \lambda_{\min}(BA) \leq \max\{\lambda_1, \lambda_2\}.$$

## 6. Numerical Tests

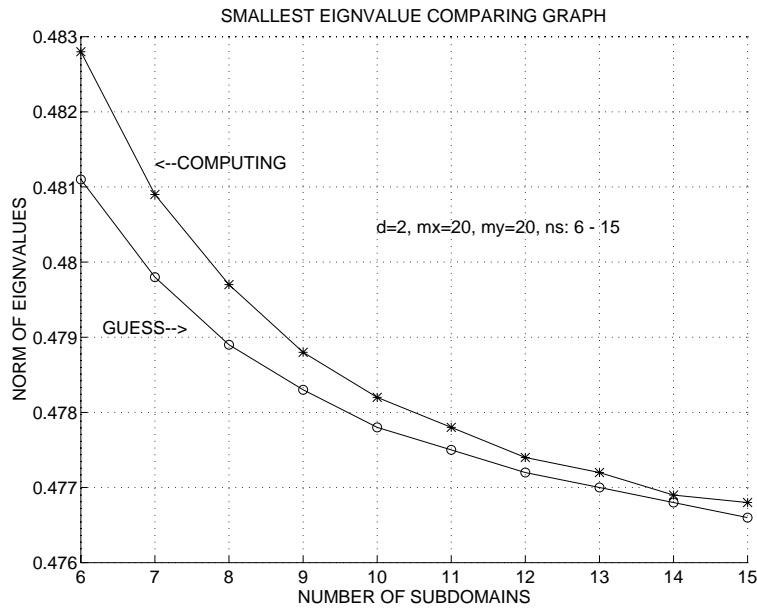
The following numerical tests have been done by Pan Feng and He Wei-jun with whom we have helpful discussion.

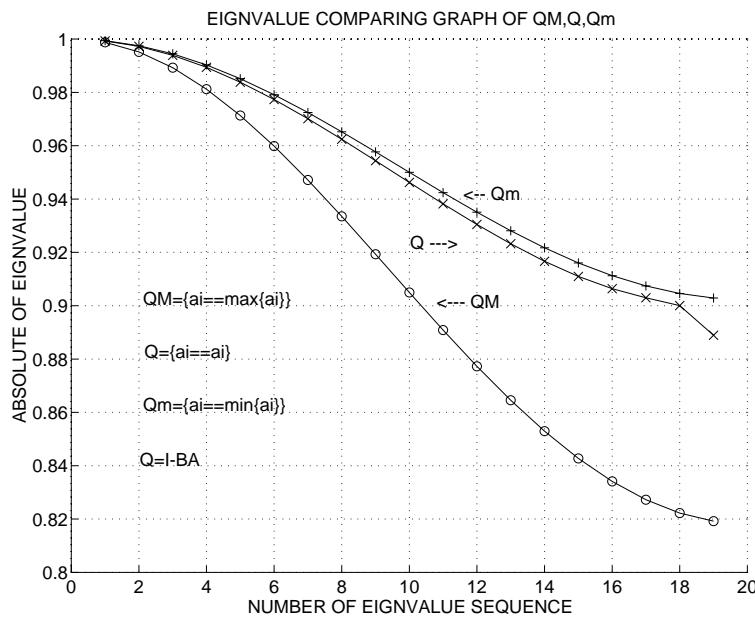
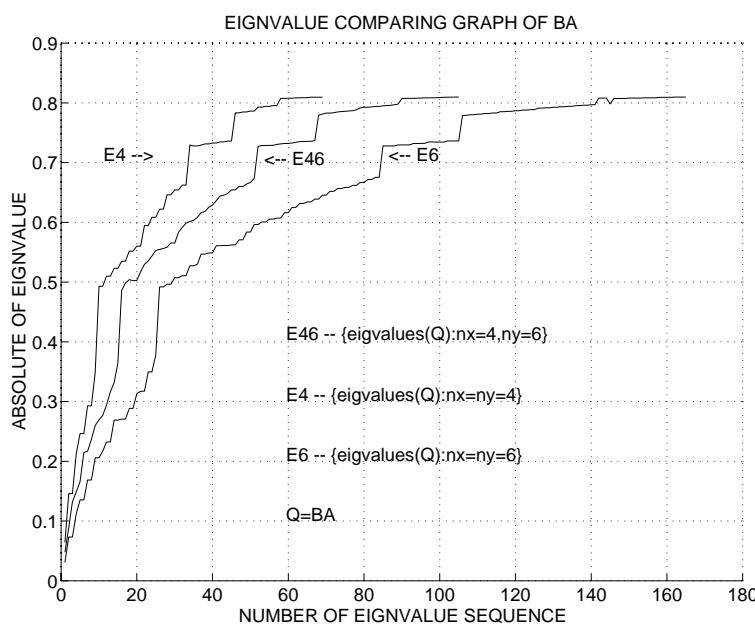
In following figures, *computing* means the smallest eigenvalue computed by MATLAB and *GUESS* means the related estimation based on our analysis.

**Test 1. 1-D model problem with uniform mesh**  $-u'' = f$ ,  $u(0) = u(1) = 0$ ,  $ns$  – number of subdomains,  $m$  – mesh number within a subdomain,  $h^{-1} = m \times ns + 1$   $d$  – overlapping mesh size = 2.



### Test 2. 1-D non-uniform case



**Test 3. 2-D Laplace equation with rectangle partition and block subdomains**

**Test 4. 2-D Laplace equation with rectangle partition and string subdomain**


## 7. Conclusion

Under a block-rectangle subdomain partitions, in this paper we have got some main results on the smallest eigenvalue estimation of  $B_s A$  as follows

$$\lambda_{min}(B_s A) \approx \frac{\pi^2}{2} H_o \hat{h}_o, \quad (14)$$

$H_o$  – Average mesh width of related 1-D subdomain,

$\hat{h}_o$  – Minimal generalized mesh width of overlapping among 1-D related subdomains.

Analogy, for a strip subdomain partitions the estimation should be modifies as

$$\lambda_{min}(B_s A) \approx \frac{\pi^3}{2} H_o \hat{h}_o, \quad (15)$$

The detailed analysis will be appeared later.

The above leading term estimations can be used for many high dimension problem decomposed from lower dimension.

It seems for strip case the leading term is  $\pi$  times bigger than for block case. However, it does not mean strip decomposition is better than block decomposition for high dimension. Because string subdomain solver is much expensive!

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