

## SMOOTHING BY CONVEX QUADRATIC PROGRAMMING <sup>\*1)</sup>

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### Abstract

In this paper, we study the relaxed smoothing problems with general closed convex constraints. It is pointed out that such problems can be converted to a convex quadratic minimization problem for which there are good programs in software libraries.

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*Key words:* Relaxed smoothing, Convex quadratic Programming.

### 1. Introduction

Let

$$x_1 < x_2 < \cdots < x_n < x_{n+1}$$

and

$$y_1, y_2, \dots, y_n, y_{n+1} = y_1.$$

The mathematical form of the problems considered in this paper is to find a twice continuous differentiable periodic function  $g(x)$  with  $g(x_{n+i}) = g(x_i)$ , such that  $g(x)$  is the optimal solution of the following problem:

$$\min \int_{x_1}^{x_{n+1}} |g''(x)|^2 dx \quad (1.1)$$

$$\text{s. t. } u \in \Omega \quad (1.2)$$

where

$$u = (u_1, u_2, \dots, u_n)^T, \quad u_i = \frac{g(x_i) - y_i}{\delta y_i}, \quad (1.3)$$

$\delta y_i, i = 1, \dots, n$  are given positive numbers and  $\Omega$  is a closed convex set in  $R^n$ . We refer the problem to *relaxed smoothing problem* whenever  $\Omega \neq \{0\}$ . For  $\Omega = \{v \in R^n \mid \|v\|_2 \leq r\}$ , the problem was investigated by Reinsch [2] and it was converted to a smooth convex unconstrained optimization. Problem (1.1) with general closed convex constraints have more applications, for example,  $\Omega = \{v \in R^n \mid \|v\|_\infty \leq r\}$  is also interesting in real problems.

It is well known that the solution of the non-relaxed problem of (1.1) is a spline function. We will prove that the solution of the relaxed smoothing problem with general closed convex constraints is the spline function  $g(x) \in C^2$  of the following form:

$$g(x) = a_i + b_i(x - x_i) + c_i(x - x_i)^2 + d_i(x - x_i)^3, \quad x \in [x_i, x_{i+1}). \quad (1.4)$$

Then the task of solving problem (1.1)-(1.2) is to find  $a_i, b_i, c_i, d_i, i = 1, \dots, n$ .

In next section, we summarize some notations and the basic relations of the spline function. Section 3 illustrates that the coefficients of the spline function can be obtained by solving a

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convex quadratic programming. Finally, in Section 4, we prove that the obtained spline function is the solution of the original problem and give our conclusions.

## 2. Notations and the Basic Relations

For analysis convenience, we need the following notations. Let  $h_i := x_{i+1} - x_i$ ,

$$D = \begin{pmatrix} \delta y_1 & & & \\ & \delta y_2 & & \\ & & \ddots & \\ & & & \delta y_n \end{pmatrix} \quad \text{and} \quad H = \begin{pmatrix} h_1 & & & \\ & h_2 & & \\ & & \ddots & \\ & & & h_n \end{pmatrix}$$

be diagonal matrices in  $R^{n \times n}$ . Denote

$$y = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix}, \quad a = \begin{pmatrix} a_1 \\ a_2 \\ \vdots \\ a_n \end{pmatrix}, \quad b = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}, \quad c = \begin{pmatrix} c_1 \\ c_2 \\ \vdots \\ c_n \end{pmatrix} \quad \text{and} \quad d = \begin{pmatrix} d_1 \\ d_2 \\ \vdots \\ d_n \end{pmatrix}.$$

Note that  $a, b, c, d$  are unknown vectors. Since  $g(x_i) = a_i$ , using these notations, the relation (1.3) can be written as

$$u = D^{-1}(a - y). \quad (2.1)$$

In addition, we need the following permutation matrix

$$P := \begin{pmatrix} 0 & 1 & & 0 \\ & \ddots & \ddots & \\ & & \ddots & 1 \\ 1 & & & 0 \end{pmatrix}.$$

For this matrix  $P$  we have  $P^T P = I$ ,

$$Pa = \begin{pmatrix} a_2 \\ \vdots \\ a_n \\ a_1 \end{pmatrix} \quad \text{and} \quad P^T a = \begin{pmatrix} a_n \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix}.$$

Now, let us list the basic properties of the periodic spline function  $g(x) \in C^2$ . First, since  $g(x_{i+1}^-) = g(x_{i+1}^+)$ , we have  $a_i + b_i h_i + c_i h_i^2 + d_i h_i^3 = a_{i+1}$  and thus

$$a + Hb + H^2c + H^3d = Pa. \quad (2.2)$$

In addition, because  $g'(x_{i+1}^-) = g'(x_{i+1}^+)$ , we have  $b_i + 2c_i h_i + 3d_i h_i^2 = b_{i+1}$  and

$$b + 2Hc + 3H^2d = Pb. \quad (2.3)$$

Finally, since  $g''(x_{i+1}^-) = g''(x_{i+1}^+)$ , we have  $c_i + 3d_i h_i = c_{i+1}$  and thus

$$c + 3Hd = Pc. \quad (2.4)$$

### 3. The Convex Quadratic Programming

If the solution of Problem (1.1)-(1.2) is a spline function of form (1.4), the objective function can be written as

$$\begin{aligned} \int_{x_1}^{x_{n+1}} |g''(x)|^2 dx &= \sum_{i=1}^n \int_{x_i}^{x_{i+1}} |2c_i + 6d_i(x - x_i)|^2 dx \\ &= \sum_{i=1}^n (4h_i c_i^2 + 12c_i d_i h_i^2 + 12d_i^2 h_i^3) \\ &= 4c^T Hc + 6c^T H^2 d + 6d^T H^2 c + 12d^T H^3 d. \end{aligned} \quad (3.1)$$

Substituting  $Hd = \frac{1}{3}(P - I)c$  (see (2.4)) in (3.1) and by a manipulation we get

$$\int_{x_1}^{x_{n+1}} |g''(x)|^2 dx = \frac{2}{3} c^T M c, \quad (3.2)$$

where

$$M = 2H + 2P^T H P + H P + P^T H. \quad (3.3)$$

Note that

$$H P = \begin{pmatrix} 0 & h_1 & & 0 \\ & \ddots & \ddots & \\ & & \ddots & h_{n-1} \\ h_n & & & 0 \end{pmatrix}, \quad P^T H P = \begin{pmatrix} h_n & & & \\ & h_1 & & \\ & & \ddots & \\ & & & h_{n-1} \end{pmatrix}$$

and thus

$$M = \begin{pmatrix} 2(h_1 + h_n) & h_1 & & & h_n \\ h_1 & 2(h_2 + h_1) & h_2 & & \\ & h_2 & \ddots & \ddots & \\ & & \ddots & \ddots & h_{n-1} \\ h_n & & & h_{n-1} & 2(h_n + h_{n-1}) \end{pmatrix}$$

is a positive definite matrix (since it is diagonal dominate).

It follows from (2.2) that

$$H^{-1}(P - I)a = b + Hc + H^2 d \quad (3.4)$$

and

$$-P^T H^{-1}(P - I)a = -P^T b - P^T Hc - P^T H^2 d. \quad (3.5)$$

From (2.4) we have

$$H^2 d = \frac{1}{3} H(P - I)c. \quad (3.6)$$

Adding (3.4) and (3.5) and using (3.6), we get

$$Qa = b - P^T b + Hc - P^T Hc + \frac{1}{3}(I - P^T)H(P - I)c, \quad (3.7)$$

where

$$Q = (I - P^T)H^{-1}(P - I). \quad (3.8)$$

It follows from (2.3) that

$$\begin{aligned} b - P^T b &= 2P^T Hc + 3P^T H^2 d \\ &\stackrel{(3.6)}{=} 2P^T Hc + P^T H(P - I)c \\ &= P^T Hc + P^T H P c. \end{aligned} \quad (3.9)$$

Substituting (3.9) into (3.7), we obtain (see (3.3))

$$Qa = \frac{1}{3}(2H + 2P^T H P + H P + P^T H)c = \frac{1}{3}M c. \quad (3.10)$$

According to (3.10), the objective function (3.2) can be rewritten as

$$6a^T Q^T M^{-1} Q a. \quad (3.11)$$

By using  $u = D^{-1}(a - y)$ , we convert the original problem to the following convex quadratic minimization problem:

$$\begin{aligned} \min \quad & \frac{1}{2}u^T D^T Q^T M^{-1} Q D u + y^T Q^T M^{-1} Q D u \\ \text{s. t} \quad & u \in \Omega. \end{aligned} \quad (3.12)$$

After getting the solution of (3.12), we can get the solution of the vectors  $a, b, c$  and  $d$  by

$$\begin{aligned} a &\stackrel{(2.1)}{=} D u + y, \\ c &\stackrel{(3.10)}{=} 3M^{-1} Q a, \\ d &\stackrel{(2.4)}{=} \frac{1}{3}H^{-1}(P c - c), \\ b &\stackrel{(2.2)}{=} H^{-1}(P a - a) - H(c + H d). \end{aligned}$$

#### 4. Optimality

The purpose of this section is to prove that the spline function (1.4) with  $a, b, c, d$  obtained from the last section is the solution of Problem (1.1)-(1.2). First, we prove the following lemma.

**Lemma 1.** *Let  $u$  be a solution of (3.12). Then we have*

$$(u' - u)^T D Q c \geq 0, \quad \forall u' \in \Omega. \quad (4.1)$$

*Proof.* Denote the objective function of (3.12) by  $\theta(u)$ . Since  $\Omega$  is closed convex and  $u$  is a solution of (3.12), it follows that  $u \in \Omega$  and

$$(u' - u)^T \nabla \theta(u) \geq 0, \quad \forall u' \in \Omega.$$

Note that

$$\nabla \theta(u) = D^T Q^T M^{-1} Q D u + D^T Q^T M^{-1} Q y.$$

Since  $D$  and  $Q$  are symmetric, it follows that

$$\begin{aligned} \nabla \theta(u) &\stackrel{(2.1)}{=} D Q M^{-1} Q (a - y) + D Q M^{-1} Q y \\ &\stackrel{(3.10)}{=} \frac{1}{3} D Q c. \end{aligned}$$

The assertion of this lemma is proved.

Now, we are in the stage to prove the optimality theorem.

**Theorem 1.** *Let  $f(x)$  be a twice continuous differentiable periodic function,  $f(x_i) = \tilde{a}_i$ ,  $f(x_i) = f(x_{n+i})$  and  $\tilde{u} = D^{-1}(\tilde{a} - y) \in \Omega$ . Then we have*

$$\int_{x_1}^{x_{n+1}} |g''(x)|^2 dx \leq \int_{x_1}^{x_{n+1}} |f''(x)|^2 dx.$$

*Proof.* Since

$$\begin{aligned} \int_{x_1}^{x_{n+1}} |f''(x)|^2 dx &= \int_{x_1}^{x_{n+1}} |g''(x)|^2 dx + \int_{x_1}^{x_{n+1}} |f''(x) - g''(x)|^2 dx \\ &\quad + 2 \int_{x_1}^{x_{n+1}} [g''(x)(f''(x) - g''(x))] dx, \end{aligned}$$

we only need to show that

$$\int_{x_1}^{x_{n+1}} g''(x)(f''(x) - g''(x)) dx \geq 0.$$

Using  $f, g \in C^2$  and by a manipulation (integration by parts), we get

$$\begin{aligned} &\int_{x_1}^{x_{n+1}} g''(x)(f''(x) - g''(x)) dx \\ &= \sum_{i=1}^n \int_{x_i}^{x_{i+1}} g''(x)(f''(x) - g''(x)) dx \\ &= \sum_{i=1}^n (f'(x) - g'(x))g''(x) \Big|_{x_i}^{x_{i+1}} - \sum_{i=1}^n \int_{x_i}^{x_{i+1}} (f'(x) - g'(x))g'''(x) dx \\ &= - \sum_{i=1}^n \int_{x_i}^{x_{i+1}} (f'(x) - g'(x))g'''(x) dx. \end{aligned} \tag{4.2}$$

The last equation of (4.2) is followed from the periodicity of  $g$ . Integrate the function again and use  $g^{(4)} = 0$ , we obtain

$$\begin{aligned} &\int_{x_1}^{x_{n+1}} g''(x)(f''(x) - g''(x)) dx \\ &= - \sum_{i=1}^n \int_{x_i}^{x_{i+1}} (f'(x) - g'(x))g'''(x) dx \\ &= - \sum_{i=1}^n (f(x) - g(x))g'''(x) \Big|_{x_i}^{x_{i+1}} + \sum_{i=1}^n \int_{x_i}^{x_{i+1}} (f(x) - g(x))g^{(4)}(x) dx \\ &= - \sum_{i=1}^n (f(x) - g(x))g'''(x) \Big|_{x_i}^{x_{i+1}} \quad (\text{since } g^{(4)} = 0) \\ &= 6 \left( (f(x_1) - a_1)(d_1 - d_n) + \sum_{i=2}^n (f(x_i) - a_i)(d_i - d_{i-1}) \right) \\ &= 6(\tilde{a} - a)^T (d - P^T d). \end{aligned} \tag{4.3}$$

Using (3.6), we obtain

$$d - P^T d = -\frac{1}{3}(P^T - I)H^{-1}(P - I)c \stackrel{(3.8)}{=} \frac{1}{3}Qc.$$

Substituting it into (4.2) and using the assertion of Lemma 1, we get

$$\begin{aligned} \int_{x_1}^{x_{n+1}} g''(x)(f''(x) - g''(x))dx &= 2(\tilde{a} - a)^T Qc \\ &= 2(D^{-1}(\tilde{a} - y) - D^{-1}(a - y))^T DQc \\ &= 2(\tilde{u} - u)^T DQc \\ &\geq 0. \end{aligned}$$

The proof is complete.

**Conclusions remark.** This paper pointed out that the relaxed smoothing problem with general closed convex constraints is equivalent to a convex quadratic programming (CQP) (3.12). For such CQP, if  $\Omega$  is a box or a polytope, many excellent numerical methods have been designed in the literature [1, 3]. Hence, it is meaningful to derive Problem (1.1)-(1.2) to a convex quadratic programming of form (3.12) for which there are good programs in software libraries.

## References

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